

Data Structure

(Java programming)

Chapter 10.



Contents

- **Dictionary**
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- *Dictionary*

- # Dictionary

Dictionary

- Like a map, a *dictionary* models a searchable collection of (*key*, *value*) entries
- Unlike a map, multiple entries with the same key are allowed
 - Cf., **put** of map vs. **insert** of dictionary
- The main operations of a dictionary are also searching, inserting and deleting entries
- Applications
 - English dictionary (same words)
 - Phone book (same names)
 - Etc.

• Dictionary

Dictionary Queue ADT

- entry `find(k)`
 - If the dictionary has an entry with key k , returns it, else, returns null
- list `findAll(k)`
 - Returns a collection (list) of all entries with key k
- entry `insert(k, v)`
 - Inserts and returns the entry (k, v)
 - Just inserts, does **NOT** replace the old entry
- entry `remove(e)`
 - Removes and returns the entry e from the dictionary, if fails, returns null
 - Cf., entry `remove(k)`: find an entry with key k and remove it
- list `entries()`
 - Returns a list of all the entries in the dictionary
- Common operations: `size()`, `isEmpty()`

• Dictionary

Dictionary Example

<i>Operation</i>	<i>Output</i>	<i>Dictionary</i>
insert(5,A)	(5,A)	(5,A)
insert(7,B)	(7,B)	(5,A), (7,B)
insert(2,C)	(2,C)	(5,A), (7,B), (2,C)
insert(8,D)	(8,D)	(5,A), (7,B), (2,C), (8,D)
insert(2,E)	(2,E)	(5,A), (7,B), (2,C), (8,D), (2,E)
find(7)	(7,B)	(5,A), (7,B), (2,C), (8,D), (2,E)
find(4)	null	(5,A), (7,B), (2,C), (8,D), (2,E)
find(2)	(2,C)	(5,A), (7,B), (2,C), (8,D), (2,E)
findAll(2)	{(2,C), (2,E)}	(5,A), (7,B), (2,C), (8,D), (2,E)
size()	5	(5,A), (7,B), (2,C), (8,D), (2,E)
remove(find(5))	(5,A)	(7,B), (2,C), (8,D), (2,E)
find(5)	null	(7,B), (2,C), (8,D), (2,E)

- **Dictionary**

Dictionary Operations

Algorithm findAll(k):

Input: A key k

Output: An iterable collection of entries with key equal to k

Create an initially-empty list L

for each entry e in $D.entries()$ **do**

if $e.getKey() = k$ **then**

$L.addLast(e)$

return L {the elements in L are the selected entries }

• Dictionary

Dictionary Operations

Algorithm insert(k, v):

Input: A key k and value v

Output: The entry (k, v) added to D

Create a new entry $e = (k, v)$

$S.addLast(e)$ $\{S \text{ is unordered}\}$

return e

Algorithm remove(e):

Input: An entry e

Output: The removed entry e

$\{ \text{We don't assume here that } e \text{ stores its location in } S \}$

for each node p in $S.nodes()$ **do**

if $p.element() = e$ **then**

$S.remove(p)$

return e

return null

$\{ \text{there is no entry } e \text{ in } D \}$

- ***Binary Search***

• Binary Search

Ordered Search Table

- An *ordered search table* (*search table* for short) is a dictionary implemented with a sorted array list
 - Stores the entries of the dictionary in an array list sorted by the keys
 - We can use an external comparator for arbitrary object keys
- Performance of ordered search tables
 - **insert** and **remove** take $O(n)$ time
 - With linear search, we need to probe n entries to find the location of the entry to be inserted or removed
 - In addition, we have to shift n entries at maximum to make room for the new entry or compact the array list after the removal
 - Whereas, we can perform find in $O(\log n)$ time using *binary search*
- A search table is effective for
 - Dictionaries on which searches are the most common operations, while insertions and removals are rarely performed
 - E.g., Credit card authorizations
 - Once a person registers the credit card information, it is retrieved every time he uses the credit card

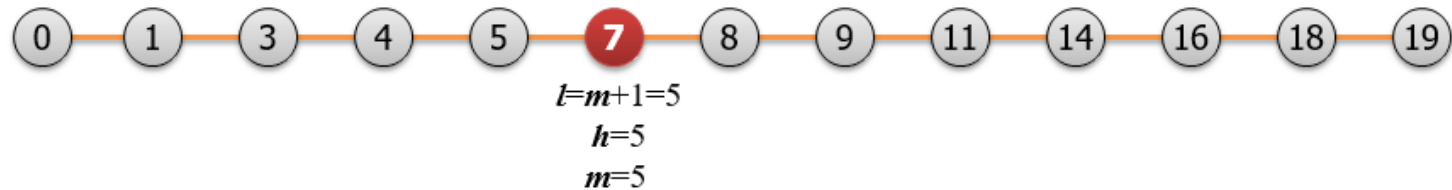
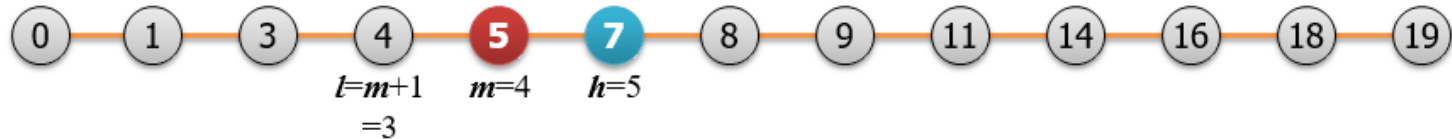
• Binary Search

Binary Search

- Example: **find(7)** from 13 integers

– Initial call: **BinarySearch(S, k=7, 0, 12)**

$n=13$



• Binary Search

Binary Search Algorithm

Algorithm BinarySearch($S, k, \text{low}, \text{high}$):

Input: An ordered array list S storing n entries and integers low and high

Output: An entry of S with key equal to k and index between low and high , if such an entry exists, and otherwise **null**

if $\text{low} > \text{high}$ **then**

return null

else

$\text{mid} \leftarrow \lfloor (\text{low} + \text{high}) / 2 \rfloor$

$e \leftarrow S.\text{get}(\text{mid})$

if $k = e.\text{getKey}()$ **then**

return e

else if $k < e.\text{getKey}()$ **then**

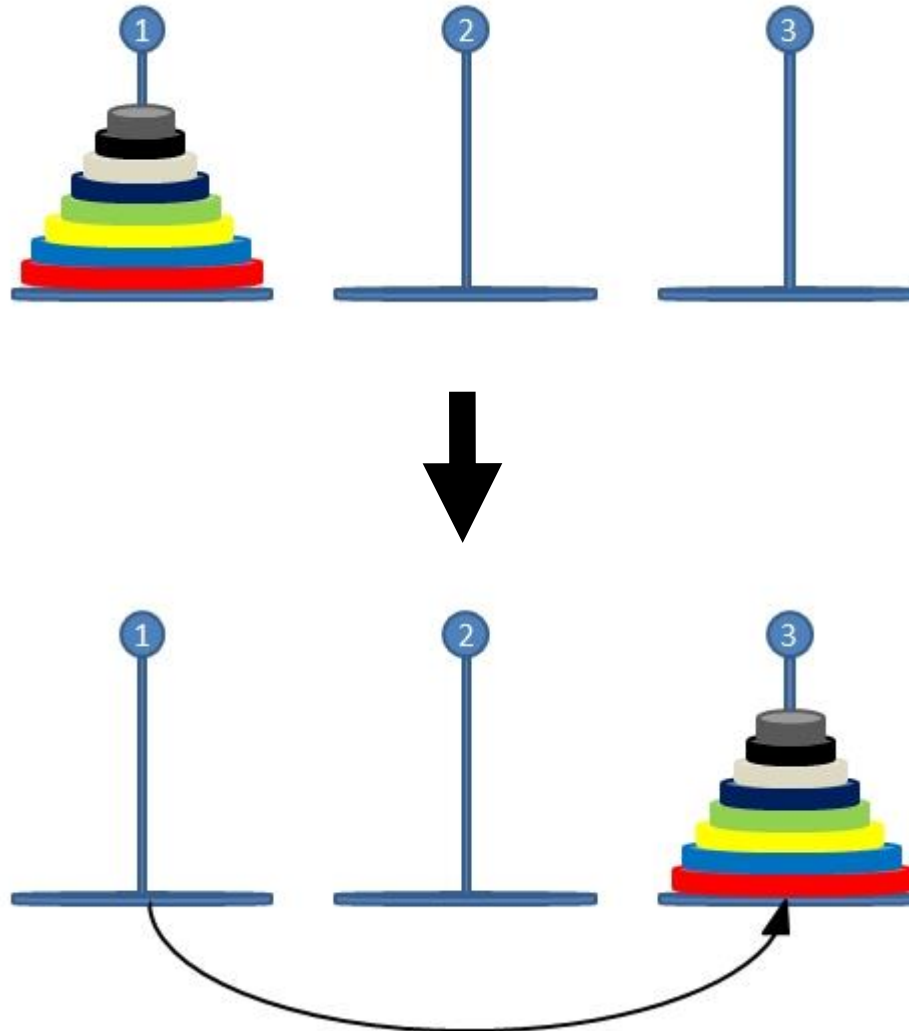
return BinarySearch($S, k, \text{low}, \text{mid} - 1$)

else

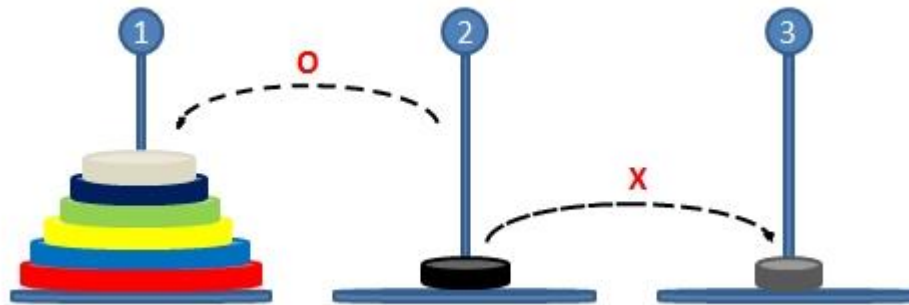
return BinarySearch($S, k, \text{mid} + 1, \text{high}$)

- *Homework3 – Tower of Hanoi with Stack*

- Tower of Hanoi with Stack



- Tower of Hanoi with Stack

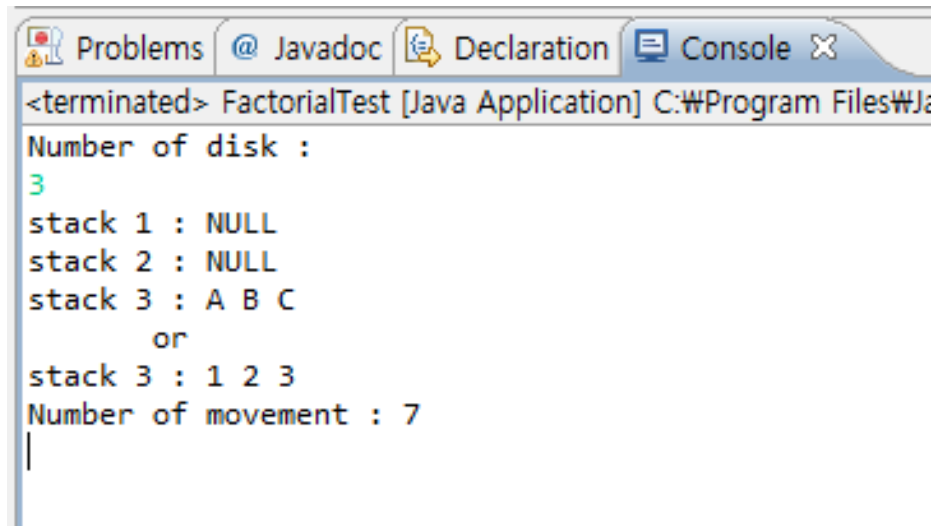


• Tower of Hanoi with Stack

- Stack을 이용하여 하노이 탑을 구현하시오
- 3개의 기둥과 각각 다른 크기의 여러 개의 원반으로 구성되어있으며, 원반은 인접해있는 기둥 사이에서만 옮길 수 있다.
- 한 기둥 안에서는 가장 큰 원반이 최하단에, 가장 작은 원반이 최상단에 위치해야 하며, 크기 순서로 쌓여야 한다.
- 처음 시작은 모든 원반이 가장 왼쪽 기둥에 쌓여있으며, 이를 모두 위의 규칙을 지키면서 가장 오른쪽 기둥으로 옮겨야 한다.
- 한 번에 하나의 원반만 옮길 수 있으며, 가장 위에 있는 원반만 꺼낼 수 있다.

• Tower of Hanoi with Stack

- Input : number of disk.
- Output : condition of three stacks, and number of movement



```
<terminated> FactorialTest [Java Application] C:\Program Files\J...
Number of disk :
3
stack 1 : NULL
stack 2 : NULL
stack 3 : A B C
      or
stack 3 : 1 2 3
Number of movement : 7
|
```

제출기한 : 2017. 05. 31 자정(24:00)

• Submission

e-mail : 2017kudatastructure@gmail.com

Attach methods :

1. Create zip file (java project folder)
2. Modify the file names :
ID_name_datastructure[Classcode].dat
ex) 201173378_이명재_datastructure_homework3[2403-1].dat
3. e-mail title :
datastructure_name_chapter
ex) datastructure_최수용_homework3[2404-1]

Classcode

2403-1 [A-1반]

2403-2 [A-2반]

2404-1 [B-1반]

2404-2 [B-2반]

2405 [C반]